Environment

- Each process has an environment, which is inherited from parent process
- Environment is a NULL-terminated array of strings
  - `extern char **environ;`
- Environment strings are of the form ‘VAR=value’
  - Variable names are capitalized by convention
getenv(3) retrieves value associated with a variable

```c
#include <stdlib.h>

char *getenv(const char *name);

char *value;

value = getenv("HOME");
if (value == NULL)
    printf("HOME not defined.\n");
else if (*value == '\0')
    printf("HOME defined but has no value.\n");
else
    printf("HOME=%s\n", value);
```
Adding Environment

- `putenv(3)` adds a var-value pair to environment
  ```c
  int putenv(const char *string)
  
  putenv("HOME=/tmp");
  ```

- `setenv(3)` also adds a var-value pair to environment
  ```c
  #include <stdlib.h>

  int setenv(const char *name, const char *value, int overwrite)
  ```
Removing environment

- `unsetenv(3)` deletes a given variable from environment

```c
#include <stdlib.h>

int unsetenv(const char *name);
```

- `clearenv(3)` clears all environment contents
getopt(3) function provides a way to handle arguments

Prototype

```c
#include <unistd.h>

int getopt(int argc, char * const argv[], const char *optstring);

extern char *optarg;
extern int optind, opterr, optopt;
```

- `optind`: index of the next argv element to be processed
- `optarg`: pointer of argument for option
Processing Arguments

Example

```c
while((opt = getopt(argc, argv, "hv:f:")) != -1)
{
    switch(opt)
    {
        case 'h':
            help();
            break;
        case 'v':
            version();
            break;
        case 'f':
            memcpy(file_name, optarg, 16);
            break;
    }
}
```
Process IDs and Process Group IDs

- **Session ID**
  - A process has a session ID
  - A session ID is set following the PID of session leader
  - Session leader = login shell
  - When a user log out every process in the session gets SIGQUIT signal

- **Process group**
  - A process belongs to a process group
  - A process group has a group leader
  - PGID = PID of group leader
  - Signals can be propagated to all processes in a group
  - This is for job controlling
  - All processes in this command belong to the same process group
    - cat ship-inventory.txt |grep booty |sort
Process IDs and Process Group IDs

Figure 5-1 illustrates the relationship between sessions, process groups, processes, and controlling terminals.

Linux provides several interfaces for setting and retrieving the session and process group associated with a given process. These are primarily of use for shells, but can also be useful to processes such as daemons that want to get out of the business of sessions and process groups altogether.

Session System Calls

Shells create new sessions on login. They do so via a special system call, which makes creating a new session easy:

```
#include <unistd.h>
pid_t setsid(void);
```

A call to `setsid()` creates a new session, assuming that the process is not already a process group leader. The calling process is made the session leader and sole member of the new session, which has no controlling tty. The call also creates a new process group inside the session and makes the calling process the process group leader and sole member. The new session’s and process group’s IDs are set to the calling process’s pid.

In other words, `setsid()` creates a new process group inside of a new session and makes the invoking process the leader of both. This is useful for daemons, which do not want to be members of existing sessions or to have controlling terminals, and for shells, which want to create a new session for each user upon login.
Process IDs and Process Group IDs

- `getpid(2)` returns PID
- `getppid(2)` returns PID of parent process

Prototype

```c
#include <sys/types.h>
#include <unistd.h>

pid_t getpid(void);
pid_t getppid(void);
```

Example

```c
#include <sys/types.h>
main()
{
    printf("My PID is %d.\n", getpid());
    printf("My PPID is %d.\n", getppid());
}
```
Process IDs and Process Group IDs

- `getpgrp(2)` returns process group ID
- `setpgid(2)` creates a new process group

Prototype

```c
int setpgrp(void);
int setpgid(pid_t pid, pid_t pgid);
```
Real and Effective IDs

- Real UID and GID can be obtained by `getuid(2)` and `getgid(2)`, respectively.
- Effective UID and GID can be obtained by `geteuid(2)` and `getegid(2)`, respectively.

Prototype

```c
#include <sys/types.h>
uid_t getuid(void);
gid_t getgid(void);
uid_t geteuid(void);
gid_t getegid(void);
```
Resource Limits

- UNIX enforces resource usage limit on each process
- Many resource limits are shown by `ulimit(1)`

```
$ ulimit -a
core file size    (blocks, -c) 0
data seg size     (kbytes, -d) unlimited
scheduling priority       (-e) 0
file size         (blocks, -f) unlimited
pending signals   (-i) 63531
max locked memory (kbytes, -l) 64
max memory size   (kbytes, -m) unlimited
open files        (-n) 1024
pipe size         (512 bytes, -p) 8
POSIX message queues (bytes, -q) 819200
real-time priority       (-r) 0
stack size        (kbytes, -s) 8192
cpu time          (seconds, -t) unlimited
max user processes (-u) 63531
virtual memory    (kbytes, -v) unlimited
file locks        (-x) unlimited
```
Resource Limits

- **Hard limit**
  - Root can lower or raise
  - Users can lower but not raise again

- **Soft limit**
  - User can lower or raise (up to hard limit)
  - Root can lower or raise

- Limits are inherited to the child processes
## Resource Limits

<table>
<thead>
<tr>
<th>Resource Macro</th>
<th>Meaning</th>
<th>Signal</th>
<th>Errno</th>
</tr>
</thead>
<tbody>
<tr>
<td>RLIMIT_CORE</td>
<td>Maximum size of a core file in bytes that may be created by a process</td>
<td></td>
<td></td>
</tr>
<tr>
<td>RLIMIT_CPU</td>
<td>Maximum amount of CPU time in seconds used by a process</td>
<td>SIGXCPU</td>
<td>EAGAIN</td>
</tr>
<tr>
<td>RLIMIT_DATA</td>
<td>Maximum size of process’s heap in bytes</td>
<td></td>
<td>ENOMEM</td>
</tr>
<tr>
<td>RLIMIT_NOFILE</td>
<td>Maximum number of open file descriptors</td>
<td></td>
<td></td>
</tr>
<tr>
<td>RLIMIT_STACK</td>
<td>Maximum size of a process’s stack in bytes</td>
<td>SIGSEGV</td>
<td></td>
</tr>
<tr>
<td>RLIMIT_NPROC</td>
<td>Maximum number of processes that can be created for a UID</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>
Resource Limits

 Prototype

```c
#include <sys/time.h>
#include <sys/resource.h>

int getrlimit(int resource, struct rlimit *rlim);
int setrlimit(int resource, const struct rlimit *rlim);

struct rlimit {
    rlim_t rlim_cur; /* Soft limit */
    rlim_t rlim_max; /* Hard limit (ceiling for rlim_cur) */
};
```
Resource Limits

Example

```c
#include <sys/resource.h>
#include <unistd.h>
main()
{
    struct rlimit myrlim;

    getrlimit(RLIMIT_NOFILE, &myrlim);
    printf("I can only open %d files\n", myrlim.rlim_cur);
    myrlim.rlim_cur = 256;
    if(setrlimit(RLIMIT_NOFILE, &myrlim) == -1)
        perror("setrlimit");
    getrlimit(RLIMIT_NOFILE, &myrlim);
    printf("I can now open %d files.\n", myrlim.rlim_cur);
    printf("sysconf() says %d files.\n", sysconf(_SC_OPEN_MAX));
}
```
You can determine the time usage of a process with `times(2)`
- Time reported by `times(2)` is in clock ticks
- You have to convert clock ticks to second

Prototype
```
#include <sys/times.h>

clock_t times(struct tms *buf);

struct tms {
    clock_t tms_utime;    /* user time */
    clock_t tms_stime;    /* system time */
    clock_t tms_cutime;   /* user time of children */
    clock_t tms_cstime;   /* system time of children */
};
```
Time Usage

- **Types of time**
  - **Wall-clock time**: time spent in real world
    - Return value of `times` shows elapsed wall-clock time from an arbitrary time point
  - **User time**: time spent in user-level
  - **System time**: time spent in kernel-level
  - **User time of children**: time spent by terminated and cleaned up children in user-level
  - **System time of children**: time spent by terminated and cleaned up children in kernel-level
# Time Usage

## Example

```c
#include <sys/types.h>
#include <sys/times.h>
#include <unistd.h>

main()
{
    int m;
    time_t t;
    struct tms mytms;
    clock_t time1, time2;
    double tick = sysconf(_SC_CLK_TCK);
    if((time1 = times(&mytms)) == -1)
    {
        perror("times"); exit(1);
    }
    for( m = 0 ; m < 99999 ; m++)
    {
        time(&t);
    }
    if((time2 = times(&mytms)) == -1)
    {
        perror("times"); exit(1);
    }
    printf("Real time: %.1f sec.\n", (time2-time1)/tick);
    printf("User time: %.1f sec.\n", mytms.tms_utime/tick);
    printf("Sys time: %.1f sec.\n", mytms.tms_stime/tick);
}
```
Current Directory

- `getcwd(3)` retrieves current working directory
- `chdir(2)` changes current working directory

Prototype

```c
#include <unistd.h>

char *getcwd(char *buf, size_t size);
int chdir(const char *path);
```
# Current Directory

## Example

```c
#include <sys/param.h>
#include <unistd.h>
main()
{
    char *dir;
    long pathmaxlen = pathconf("/", _PC_PATH_MAX);
    dir=getcwd((char *)NULL, pathmaxlen+1);
    if(dir==NULL)
        {perror("getcwd"); exit(1)}
    printf("CWD: %s\n", dir);
    free(dir);
    if(chdir("/tmp") == -1)
        perror("chdir");
    dir=getcwd((char *)NULL, pathmaxlen+1);
    if(dir==NULL)
        perror("getcwd");
    printf("CWD: %s\n", dir);
}
```
Shell Lab

- **Skeleton of a shell is very simple**

  ```
  while(1) {
    print prompt
    read command
    process command
  }
  ```

- **Let’s make a shell!**