

# Representing and Manipulating Integers Part I

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# Unsigned Integers

- Encoding unsigned integers

$$B = [b_{w-1}, b_{w-2}, \dots, b_0] \quad \mathbf{x} = 0000\ 0111\ 1101\ 0011_2$$

$$D(B) = \sum_{i=0}^{w-1} b_i \cdot 2^i$$

$$\begin{aligned} D(\mathbf{x}) &= 2^{10} + 2^9 + 2^8 + 2^7 \\ &\quad + 2^6 + 2^4 + 2^1 + 2^0 \\ &= 1024 + 512 + 256 + 128 \\ &\quad + 64 + 16 + 2 + 1 \\ &= 2003 \end{aligned}$$

- What is the range for unsigned values with  $w$  bits?

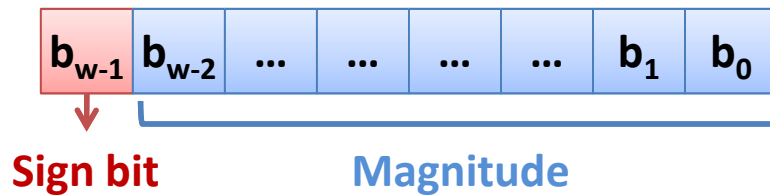
# Signed Integers (1)



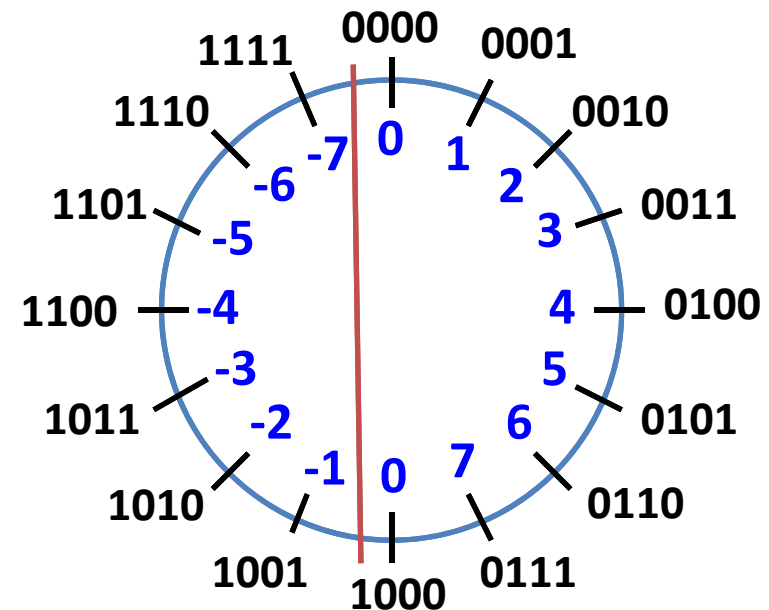
- **Encoding positive numbers**
  - Same as unsigned numbers
  
- **Encoding negative numbers**
  - Sign-magnitude representation
  - Ones' complement representation
  - Two's complement representation

# Signed Integers (2)

## ■ Sign-magnitude representation



$$S(B) = (-1)^{b_{w-1}} \cdot \left( \sum_{i=0}^{w-2} b_i \cdot 2^i \right)$$



- Two zeros
  - [000...00], [100...00]
- Used for floating-point numbers

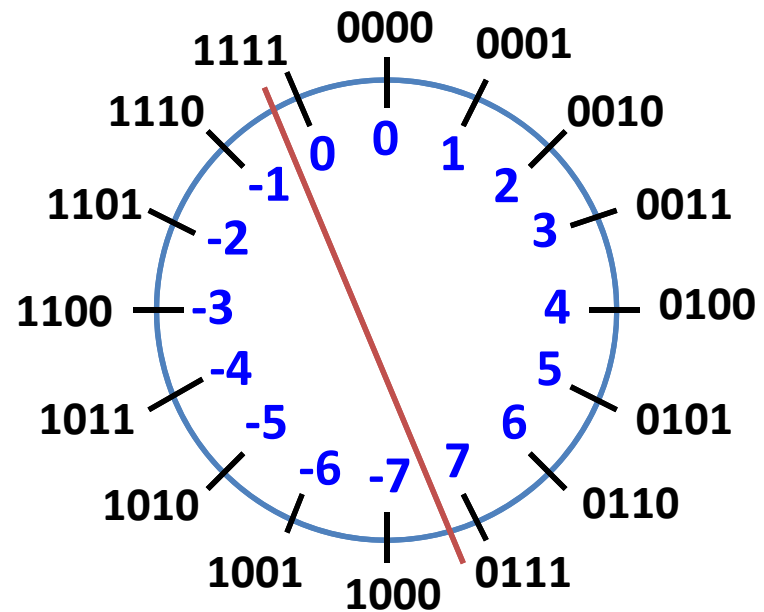
# Signed Integers (3)

## ▪ Ones' complement representation



Sign bit

$$O(B) = -b_{w-1}(2^{w-1} - 1) + \left( \sum_{i=0}^{w-2} b_i \cdot 2^i \right)$$



- Easy to find  $-n$
- Two zeros
  - $[000\dots00]$ ,  $[111\dots11]$
- No longer used

# Signed Integers (4)

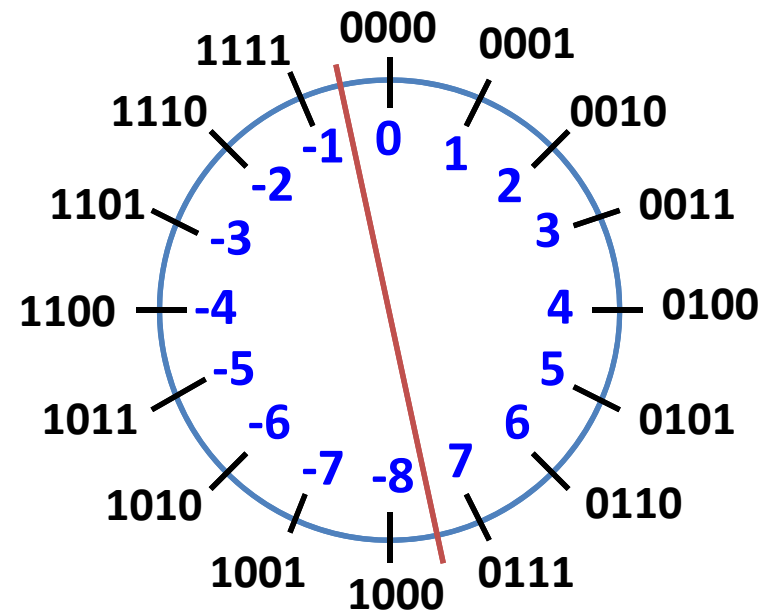
## Two's complement representation



Sign bit

$$O(B) = -b_{w-1} \cdot 2^{w-1} + \left( \sum_{i=0}^{w-2} b_i \cdot 2^i \right)$$

- Unique zero
- Easy for hardware
  - leading 0  $\geq 0$
  - leading 1  $< 0$
- Used by almost all modern machines



# Signed Integers (5)

## ■ Two's complement representation (cont'd)

- Following holds for two's complement

$$\sim x + 1 == -x$$

- Complement

– Observation:  $\sim x + x == 1111\dots11_2 == -1$

- Increment

$$\sim x + x == -1$$

$$\sim x + x + (-x + 1) == -1 + (-x + 1)$$

$$\sim x + 1 == -x$$

# Numeric Ranges (1)

## ■ Unsigned values

- UMin = 0 [000...00]
- UMax =  $2^w - 1$  [111...11]

## ■ Two's complement values

- TMin =  $-2^{w-1}$  [100...00]
- TMax =  $2^{w-1} - 1$  [011...11]

Values for  $w = 16$

	Decimal	Hex	Binary
UMax	65535	FF FF	11111111 11111111
TMax	32767	7F FF	01111111 11111111
TMin	-32768	80 00	10000000 00000000
-1	-1	FF FF	11111111 11111111
0	0	00 00	00000000 00000000



# Numeric Ranges (2)

## ■ Values for different word sizes

	w = 8	w = 16	w = 32	w = 64
UMax	255	65,535	4,294,967,295	18,446,744,073,709,551,615
TMax	127	32,767	2,147,483,647	9,223,372,036,854,775,807
TMin	-128	-32,768	-2,147,483,648	-9,223,372,036,854,775,808

## ■ Observations

- $|TMin| = TMax + 1$   
(Asymmetric range)
- $UMax = 2 * TMax + 1$

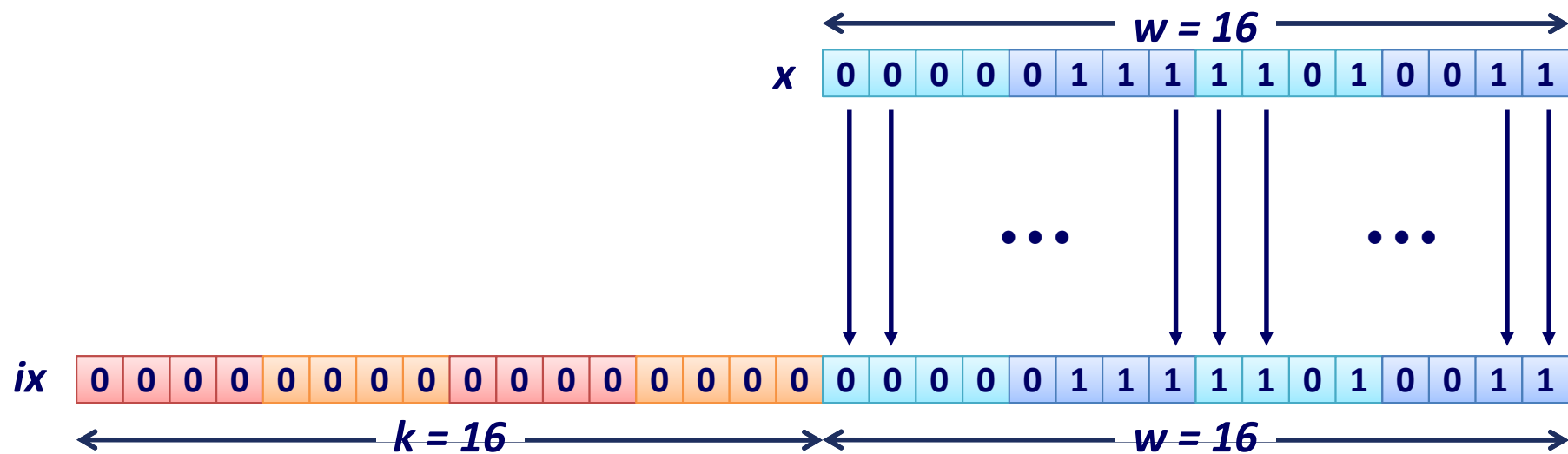
## ■ In C programming

- `#include <limits.h>`
- `INT_MIN`, `INT_MAX`,  
`LONG_MIN`, `LONG_MAX`,  
`UINT_MAX`, ...
- Values platform-specific

# Type Conversion (1)

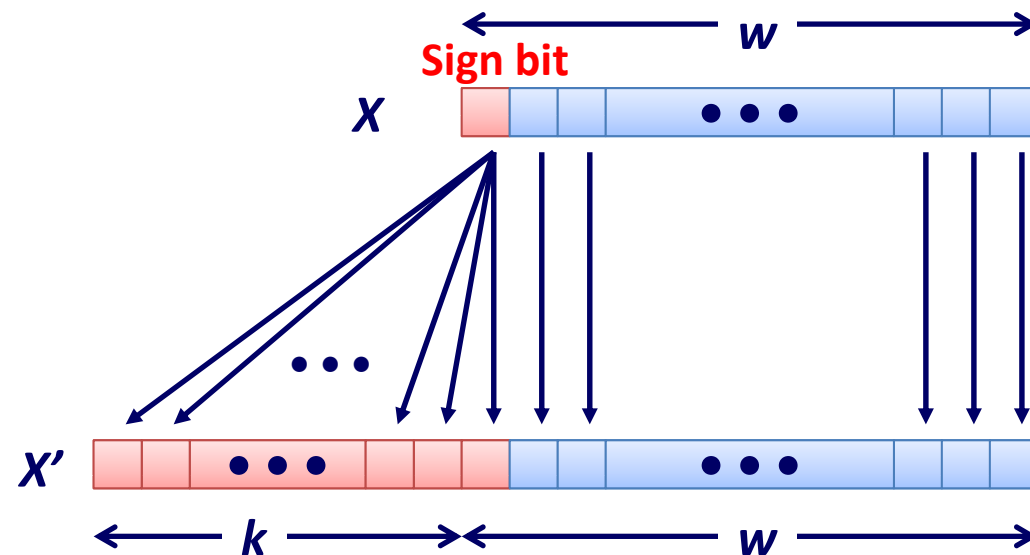
- Unsigned:  $w$  bits  $\rightarrow w+k$  bits
  - Zero extension: just fill  $k$  bits with 0's

```
unsigned short x = 2003;
unsigned       ix = (unsigned) x;
```



# Type Conversion (2)

- **Signed:**  $w$  bits  $\rightarrow w+k$  bits
  - Given  $w$ -bit signed integer  $x$
  - Convert it to  $w+k$ -bit integer with same value
- **Sign extension**
  - Make  $k$  copies of sign bit



# Type Conversion (3)

## ▪ Sign extension example

- Converting from smaller to larger integer type
- C automatically performs sign extension

```
short int x = 2003;
int      ix = (int) x;
short int y = -2003;
int      iy = (int) y;
```

	Decimal	Hex	Binary
x	2003	07 D3	00000111 11010011
ix	2003	00 00 07 D3	00000000 00000000 00000111 11010011
y	-2003	F8 2D	11111000 00101101
iy	-2003	FF FF F8 2D	11111111 11111111 11111000 00101101

# Type Conversion (4)

- **Unsigned & Signed:  $w+k$  bits  $\rightarrow$   $w$  bits**
  - Just truncate it to lower  $w$  bits
  - Equivalent to computing  $x \bmod 2^w$

```
unsigned int    x = 0xcafebabe;
unsigned short ix = (unsigned short) x;
int            y = 0x2003beef;
short         iy = (short) y;
```

	Decimal	Hex	Binary
x	3405691582	CA FE BA BE	11001010 11111110 10111010 10111110
ix	47806	BA BE	10111010 10111110
y	537116399	20 03 BE EF	00100000 00000011 10111110 11101111
iy	-16657	BE EF	10111110 11101111

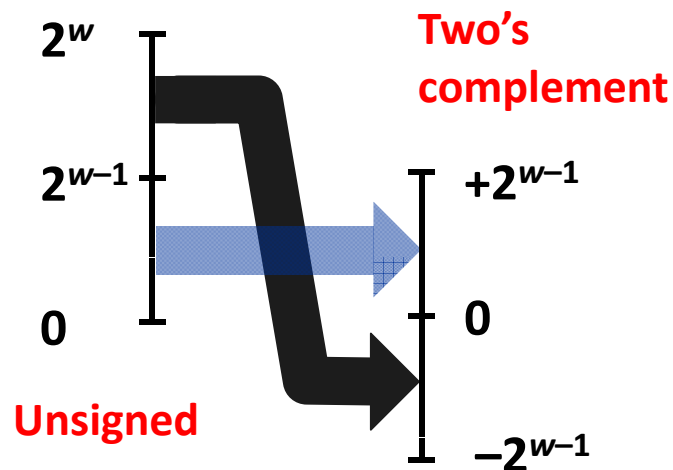
# Type Conversion (5)

## ■ Unsigned → Signed

- The same bit pattern is interpreted as a signed number

$$U2T_w(x) = \begin{cases} x, & x < 2^{w-1} \\ x - 2^w, & x \geq 2^{w-1} \end{cases}$$

```
unsigned short x = 2003;
short          ix = (short) x;
unsigned short y = 0xbabe;
short          iy = (short) y;
```



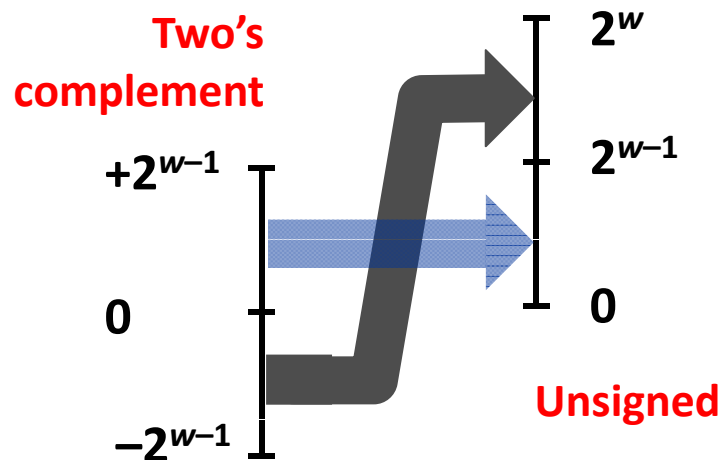
	Decimal	Hex	Binary
x	2003	07 D3	00000111 11010011
ix	2003	07 D3	00000111 11010011
y	47806	BA BE	10111010 10111110
iy	-17730	BA BE	10111010 10111110

# Type Conversion (6)

## ▪ Signed $\rightarrow$ Unsigned

- Ordering inversion
- Negative  $\rightarrow$  Big positive

$$T2U_w(x) = \begin{cases} x + 2^w, & x < 0 \\ x, & x \geq 0 \end{cases}$$



```
short      x = 2003;
unsigned short ix = (unsigned short) x;
short      y = -2003;
unsigned short iy = (unsigned short) y;
```

	Decimal	Hex	Binary
x	2003	07 D3	00000111 11010011
ix	2003	07 D3	00000111 11010011
y	-2003	F8 2D	11111000 00101101
iy	63533	F8 2D	11111000 00101101

# Type Casting in C (1)

## ■ Constants

- By default, considered to be signed integers
- Unsigned if have "U" or "u" as suffix
  - 0U, 12345U, 0x1A2Bu

## ■ Type casting

- Explicit casting
- Implicit casting via
  - Assignments
  - Procedure calls

```
int      tx, ty;
unsigned ux, uy;
tx = (int) ux;
uy = (unsigned) ty;
```

```
int f(unsigned);
tx = ux;
f(ty);
```



# Type Casting in C (2)

## ■ Expression evaluation

- If mix unsigned and signed in single expression, signed values implicitly cast to **unsigned**.
- Including comparison operations `<`, `>`, `==`, `<=`, `>=`

Expression	Type	Evaluation
<code>0 == 0U</code>	unsigned	True
<code>-1 &lt; 0</code>	signed	True
<code>-1 &lt; 0U</code>	unsigned	?
<code>-1 &gt; -2</code>	signed	?
<code>(unsigned) -1 &gt; -2</code>	unsigned	?
<code>2147483647 &gt; -2147483647-1</code>	signed	?
<code>2147483647U &gt; -2147483647-1</code>	unsigned	?
<code>2147483647 &gt; (int) 2147483648U</code>	signed	?

# Type Casting in C (3)

- **Example 1-1**

```
int main ()
{
    unsigned i;
    for (i = 10; i > 0; i--)
        printf ("%u\n", i);
}
```

- **Example 1-2**

```
int main ()
{
    unsigned i;
    for (i = 10; i >= 0; i--)
        printf ("%u\n", i);
}
```

# Type Casting in C (4)

- Example 2

```
int sum_array (int a[], unsigned len)
{
    int i;
    int result = 0;

    for (i = 0; i <= len - 1; i++)
        result += a[i];

    return result;
}
```

# Type Casting in C (5)

- Example 3-1

```
void copy_mem1 (char *src, char *dest, unsigned len)
{
    unsigned i;
    for (i = 0; i < len; i++)
        *dest++ = *src++;
}
```

- Example 3-2

```
void copy_mem2 (char *src, char *dest, unsigned len)
{
    int i;
    for (i = 0; i < len; i++)
        *dest++ = *src++;
}
```

# Type Casting in C (6)

- Example 3-3

```
void copy_mem3 (char *src, char *dest, unsigned len)
{
    for (; len > 0; len--)
        *dest++ = *src++;
}
```

- Example 3-4

```
void copy_mem4 (char *src, char *dest, unsigned len)
{
    for (; (int) len > 0; len--)
        *dest++ = *src++;
}
```

# Type Casting in C (7)

- Example 4

```
#include <stdio.h>

int main ()
{
    unsigned char    c;

    while ((c = getchar()) != EOF)
        putchar (c);
}
```

# Type Casting in C (8)

## ■ Lessons

- There are many tricky situations when you use unsigned integers – hard to debug
- Do not use just because numbers are nonnegative
- Use only when you need collections of bits with no numeric interpretation (“flags”)
- Few languages other than C support unsigned integers