

# Concurrent Programming

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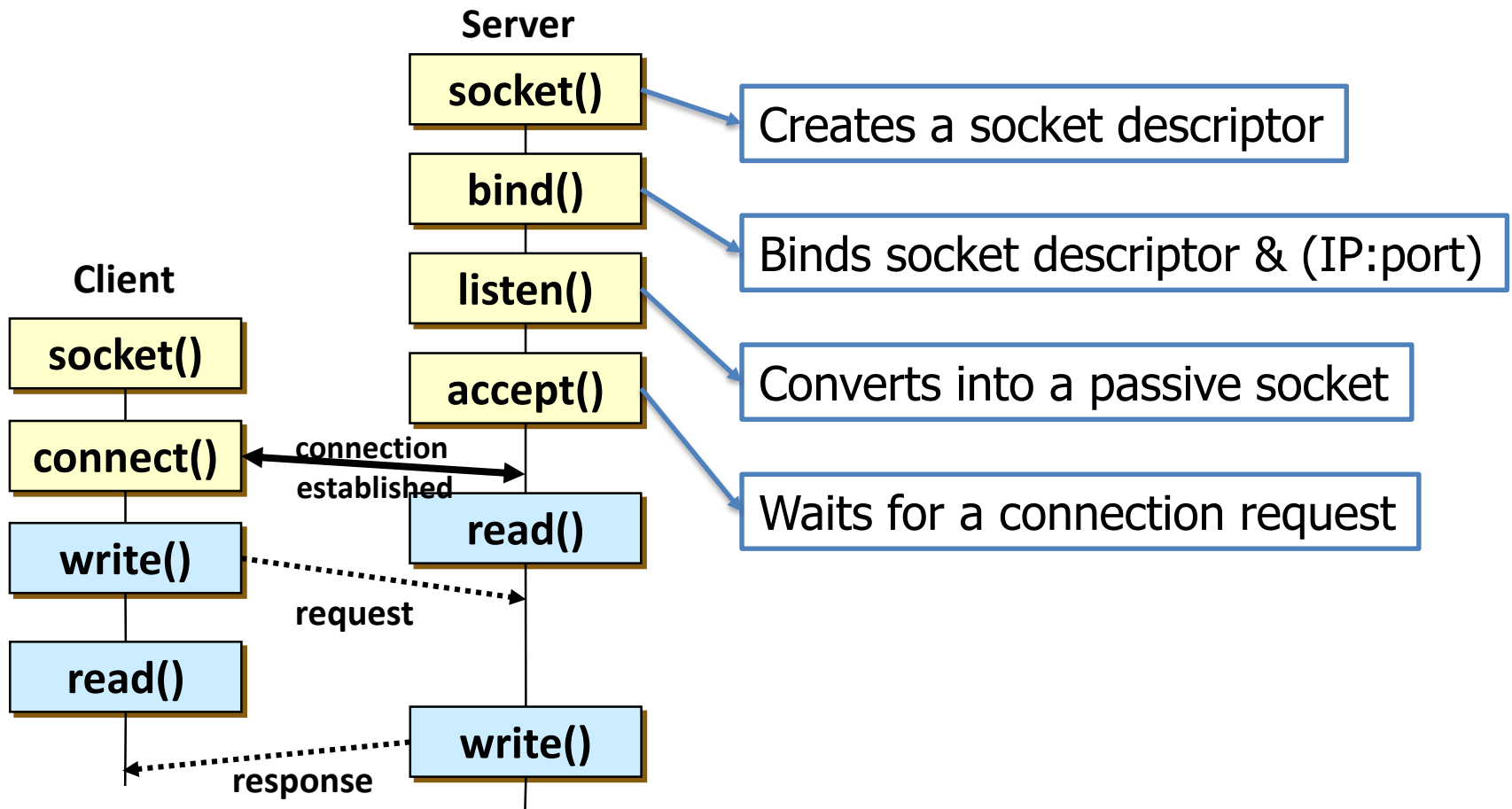
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# Server

## Connection-oriented service



# Echo Server Revisited

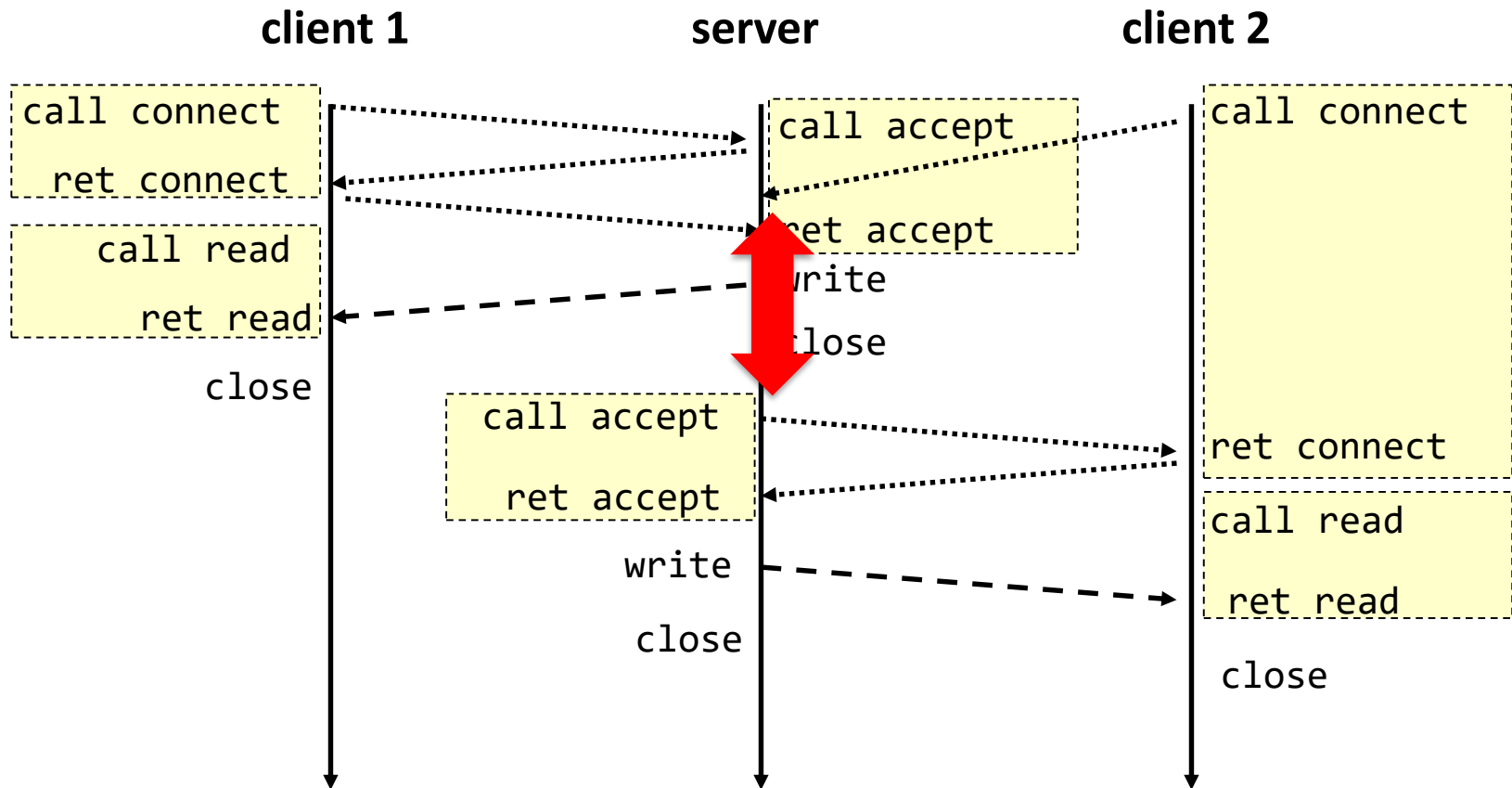
```
int main (int argc, char *argv[]) {
    ...
    listenfd = socket(AF_INET, SOCK_STREAM, 0);

    bzero((char *)&saddr, sizeof(saddr));
    saddr.sin_family = AF_INET;
    saddr.sin_addr.s_addr = htonl(INADDR_ANY);
    saddr.sin_port = htons(port);
    bind(listenfd, (struct sockaddr *)&saddr, sizeof(saddr));

    listen(listenfd, 5);
    while (1) {
        connfd = accept(listenfd, (struct sockaddr *)&caddr, &crlen);
        while ((n = read(connfd, buf, MAXLINE)) > 0) {
            printf ("got %d bytes from client.\n", n);
            write(connfd, buf, n);
        }
        close(connfd);
    }
}
```

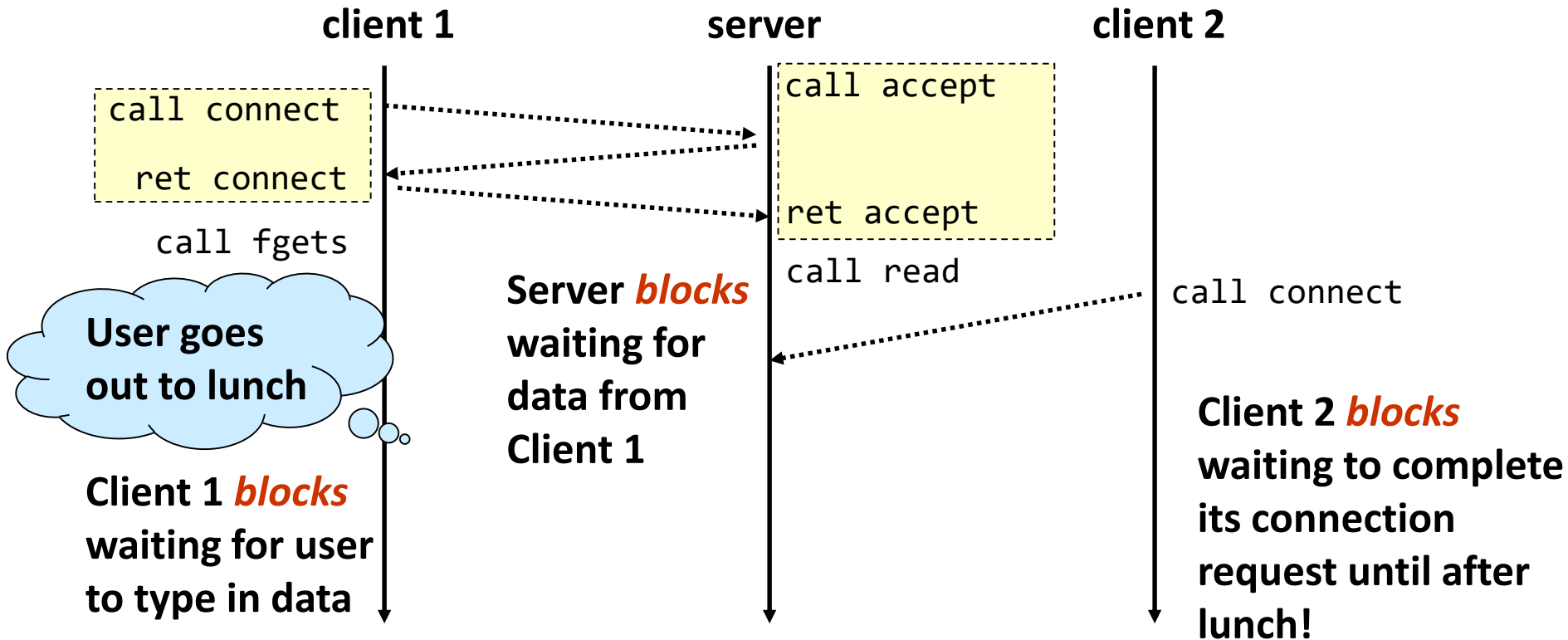
# Iterative Servers (1)

- One request at a time



# Iterative Servers (2)

## ▪ Fundamental flaw



## ▪ Solution: use concurrent servers instead

- Use multiple concurrent flows to serve multiple clients at the same time.

# Creating Concurrent Flows

## ■ Processes

- Kernel automatically interleaves multiple logical flows.
- Each flow has its own private address space.

## ■ Threads

- Kernel automatically interleaves multiple logical flows.
- Each flow shares the same address space.
- Hybrid of processes and I/O multiplexing

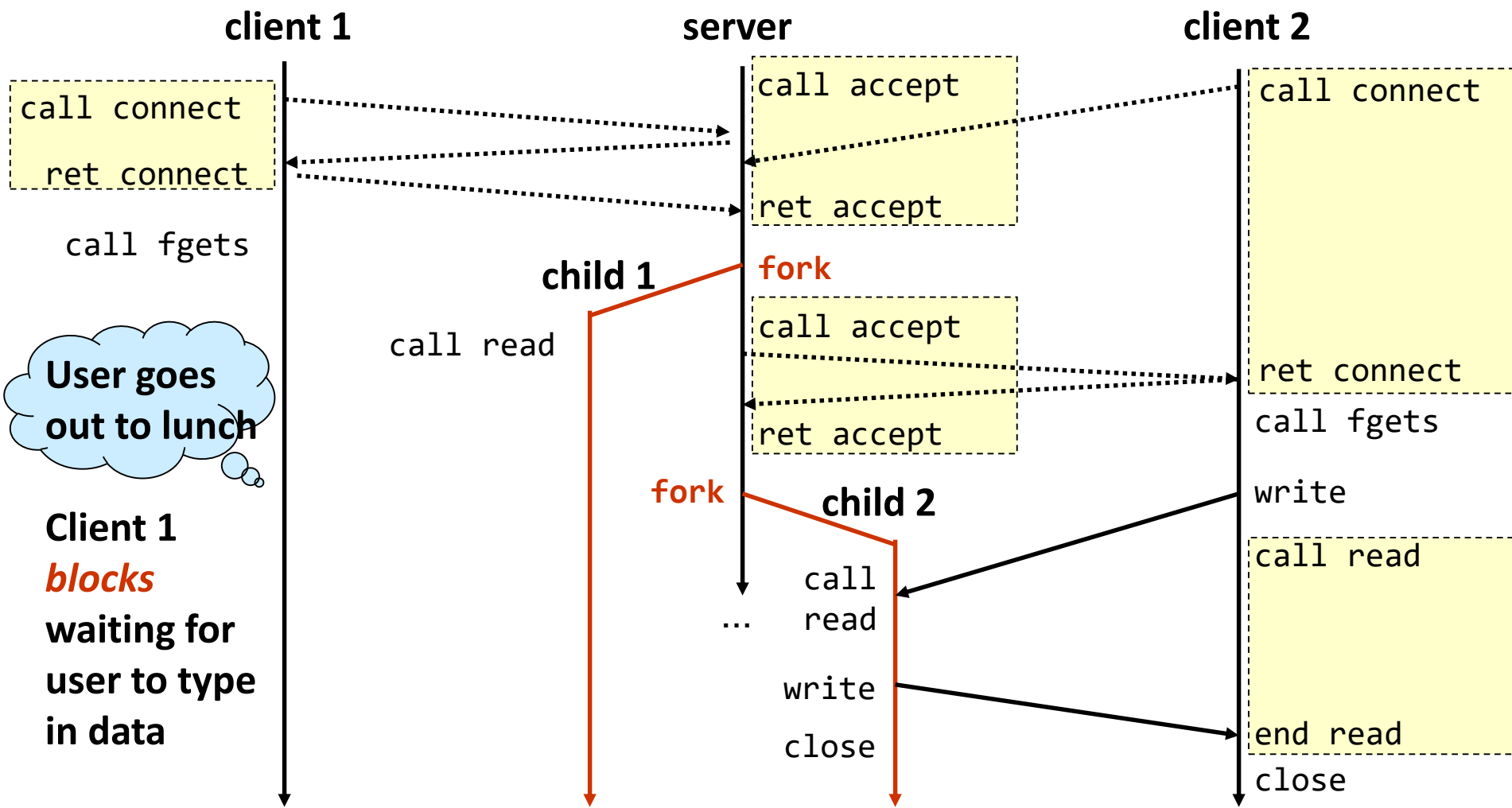
## ■ I/O multiplexing with `select()`

- User manually interleaves multiple logical flows
- Each flow shares the same address space
- Popular for high-performance server designs.

**Concurrent Programming**

**Process-based**

# Process-based Servers





# Implementation Issues

- **Servers should restart `accept()` if it is interrupted by a transfer of control to the `SIGCHLD` handler**
  - Not necessary for systems with POSIX signal handling.
  - Required for portability on some older Unix systems.
- **Server must reap zombie children**
  - to avoid fatal memory leak
- **Server must close its copy of `connfd`.**
  - Kernel keeps reference for each socket.
  - After `fork()`, `refcnt(connfd) = 2`
  - Connection will not be closed until `refcnt(connfd) = 0`

# Process-based Designs



## ■ Pros

- Handles multiple connections concurrently.
- Clean sharing model.
  - Descriptors (no), file tables (yes), global variables (no)
- Simple and straightforward.

## ■ Cons

- Additional overhead for process control.
  - Process creation and termination
  - Process switching
- Nontrivial to share data between processes.
  - Requires IPC (InterProcess Communication) mechanisms: FIFO's, System V shared memory and semaphores

# Exercise

## ■ "Guess the Number"

- At the same time, multiple client can be served by echo server
- **There should be no memory leakage**
  - How about closing files?

```
./server [port]
```

```
./client 127.0.0.1 [port]
Guess? 32
Up
Guess? 74
Down
Guess? 34
Correct!
$
```

```
./client 127.0.0.1 [port]
Guess? 34
Up
Guess? 40
Down
Guess? 35
Correct!
$
```

```
./client 127.0.0.1 [port]
Guess? 35
Down
Guess? 10
Up
Guess? 14
Correct!
$
```