

SWE3004: Operating Systems
prof. Euseong Seo

Project 3. Synchronization

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TAs

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Project plan

- Total 7 projects

- 0) Starting xv6 operating system (5%)

- 1) System call (10%)

- 2) Thread (15%)

- 3) Synchronization (15%)**

- 4) Scheduling 1(10%)

- 5) Scheduling 2(15%)

- 6) Page fault handler (15%)

- 7) Copy on Write (15%)

Locks of xv6

- **spinlock** : busy wait until lock unlocks (spinlock.c)
 - `acquire(struct spinlock *lock)`
 - `release(struct spinlock *lock)`

- **sleeplock** : sleep until lock unlocks (sleeplock.c)
 - `acquiresleep(struct sleeplock *lock)`
 - `releasesleep(struct sleeplock *lock)`
 - You don't have to use sleeplock for project3

Project 3. Mutex & CV

- Implement 6 new system calls
 - `int mutex_init(mutex_t *mutex)`
 - `int mutex_lock(mutex_t *mutex)`
 - `int mutex_unlock(mutex_t *mutex)`
 - `int cond_init(cond_t *cond)`
 - `int cond_wait(cond_t *cond, mutex_t *mutex)`
 - `int cond_signal(cond_t *cond)`

Project 3. Mutex & CV

- mutex & cv structure
 - make your own `synch.c` and `synch.h` files
 - Example :

```
#define NTHREAD 8
struct mutex_t{
    int valid;
    struct spinlock lock;
    struct proc *current;
    struct proc *queue[NTHREAD-1];
    int qsize;
};

struct cond_t{
    int active;
    struct spinlock lock;
    struct proc *queue[NTHREAD-1];
    int qsize;
};
```

valid : is initialized?
lock : spinlock for mutex
current : current locked thread
queue : queueing thread list
qsize : queue size

You don't have to use this structure. This is just sample.

mutex_init()

- Synopsis

- `int mutex_init(struct mutex_t *mutex);`

- Description

- The `mutex_init()` function initializes the mutex referenced by `mutex`. Upon successful initialization, the state of the mutex becomes initialized and unlocked.

- Return value

- return 0 : at initialization success
 - return -1 : when the mutex is invalid
 - return -2 : when attempting to reinitialize an already initialized mutex

mutex_lock()

- Synopsis

- `int mutex_lock(struct mutex_t *mutex);`

- Description

- The mutex object referenced by `mutex` is locked by calling `mutex_lock()`. If the mutex is already locked, the calling thread blocks until the mutex becomes available.

- Return value

- return 0 : when lock is achieved.
 - return -1 : when the mutex is invalid
 - return -2 : when the mutex is not initialized
 - return -3 : when calling thread already has the mutex

mutex_unlock()

- Synopsis

- `int mutex_unlock(struct mutex_t *mutex);`

- Description

- The `mutex_unlock()` function releases the mutex object referenced by `mutex`. If there are threads blocked on the mutex, the thread waiting for the mutex should be unblocked **in order** and put on the list of ready threads.

- Return value

- return 0 : when lock released successfully

- return -1 : when the mutex is invalid

- return -2 : when the mutex is not initialized

- return -3 : when calling thread does not have the mutex

cond_init()

- Synopsis

- `int cond_init(struct cond_t *cond);`

- Description

- The `cond_init()` function initializes the condition variable referenced by `cond`. Upon successful initialization, the state of the condition variable becomes initialized.

- Return value

- return 0 : at initialization success
 - return -1 : when the condition variable is invalid
 - return -2 : when attempting to reinitialize an already initialized condition variable

cond_wait()

- Synopsis

- `int cond_wait(struct cond_t *cond, struct mutex_t *mutex);`

- Description

- The `cond_wait()` function blocks on a condition variable. It should be called with mutex locked by the calling thread. It should release mutex and cause the calling thread to block on the condition variable `cond`. Upon successful return, the mutex should be locked and owned by the calling thread.

- Return value

- return 0 : at success.
 - return -1 : when the condition variable is invalid
 - return -2 : when the condition variable or mutex is not initialized
 - return -3 : when calling thread does not have the mutex

cond_signal()

- Synopsis

- `int cond_signal(struct cond_t *cond);`

- Description

- The `cond_signal()` function unblocks a thread blocked on the specified condition variable `cond`. If more than one thread is blocked on the condition variable, thread waiting for the condition variable should be unblocked **in order** and put on the list of ready threads.

- Return value

- return 0 : at success.

- return -1 : when the condition variable is invalid

- return -2 : when the condition variable is not initialized

CV semantic

- Mesa semantic
 - `signal()` places a waiter on the ready queue, but signaler continues inside the critical section.
 - You have to use Mesa semantic for this project.

Submission

- You should enter your ID & project no. (this time is 3) on **Makefile**
- `$make tarball`
- Then, `xv6-project-3-studentID.tar.gz` will be created in the parent folder.
- You need to submit a **document**.

Submission

- Send your code file (xv6-project-3-studentID.tar.gz) and document file to gmj03003@gmail.com
- Please send a mail with tittle including [SWE3004-P3]
 - Ex) [SWE3004-P3] 2014111111-project3
- **PLEASE DO NOT COPY**
 - **YOU WILL GET F GRADE IF YOU COPIED**
- Due date: 5/1(Tue.), 23:59:59 PM
 - Delays are allowed only one week from the deadline. And there will be up to -40%(6%*6 day + 4%*1 last day) penalty.

Questions

- If you have questions, please email to TA
- You can also visit #85533. Please email TA before visiting